

# Hot Potatoes Heat Up BGP Routing

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with

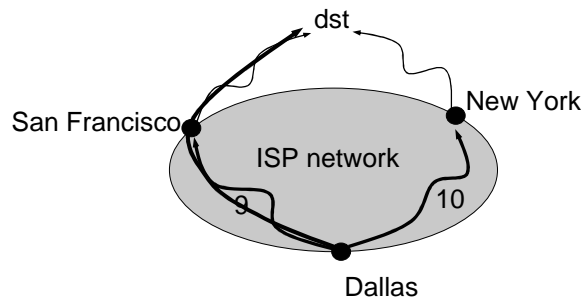
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*30<sup>th</sup> NANOG – Miami, Florida*

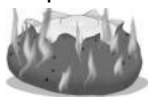
## BGP Routing inside an AS

- ♦ Most people think...
  - BGP = inter-domain routing
  - Routing inside an AS = OSPF/IS-IS
- ♦ But in large transit ASes...
  - Most traffic is routed using BGP
  - IGP changes affect BGP decisions

# Hot-Potato Routing



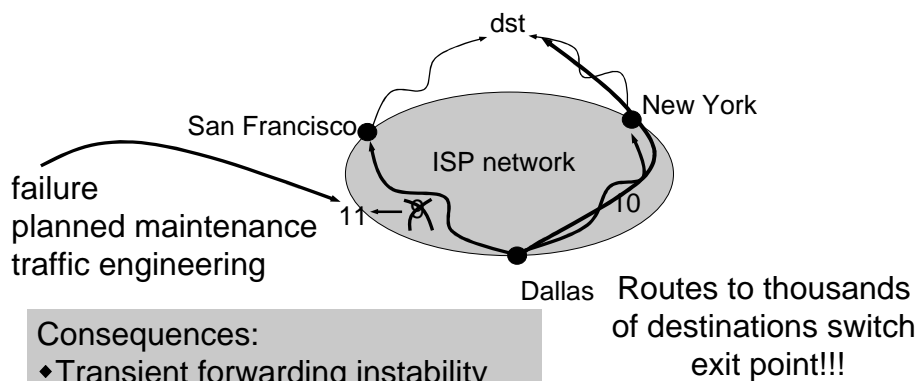
Hot-potato routing = route to closest exit point  
when there is more than one  
route to destination



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# Hot-Potato Routing




Consequences:

- ◆ Transient forwarding instability
- ◆ Traffic shift
- ◆ Inter-domain routing changes

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## BGP Decision Process

- “equally good”
- ◆ **Ignore if exit point unreachable**
  - ◆ Highest local preference
  - ◆ Lowest AS path length
  - ◆ Lowest origin type
  - ◆ Lowest MED (with same next hop AS)
  - ◆ **Lowest IGP cost to next hop** ← 
  - ◆ Lowest router ID of BGP speaker

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## Why Care about Hot Potatoes?

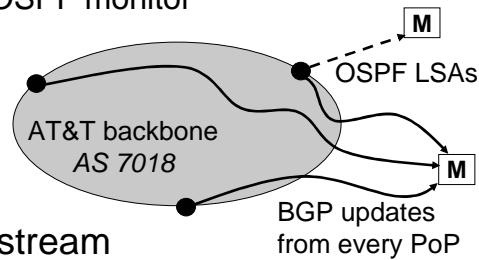
- ◆ Large number of routes potentially affected
  - All routes from peers and multi-homed customers
- ◆ Understanding of impact in real networks
  - How often hot-potato changes happen in a real network and how many prefixes do they affect?
  - What are the convergence delays?
- ◆ Avoiding routing instability
  - Operators: avoid hot-potato changes
  - Vendors: reduce impact of hot-potato changes

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# Measuring Hot-Potato Routing

- ◆ Collect measurement of both protocols
  - BGP monitor and OSPF monitor



- ◆ Pre-process each stream
  - Clustering related messages
- ◆ Correlate the two streams of data
  - Match BGP updates with OSPF events

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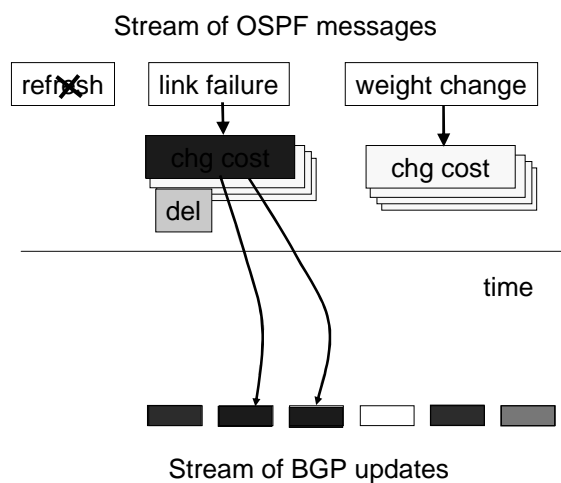
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# Heuristic for Matching

Transform stream of OSPF messages into routing changes

Match BGP updates with OSPF events that happen close in time

Classify BGP updates by possible OSPF causes



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## Impact of OSPF/BGP Interaction (June 2003)

- ♦ High variability according to location and day
  - Impact on external BGP measurements and customers

location	% updates		days > 10% of hot-potato changes
	min	max	
rich peering	0%	5.3%	0
no peering	0%	46.5%	5

- ♦ One LSA can have a big impact

location	LSAs with no impact	prefixes impacted by an LSA
rich peering	95.1%	less than 1%
no peering	94.8%	59%

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## Delay for BGP Routing Change

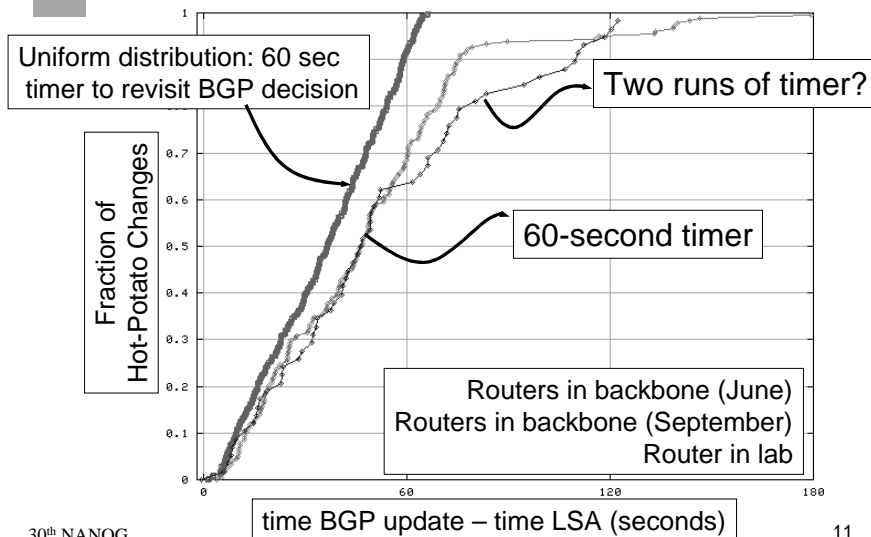
- ♦ Router rerunning BGP decision process
  - Implementation specific
    - Timer driven (tunable parameter)
    - Event driven
- ♦ Internal BGP route propagation delay
  - Internal BGP hierarchy (route reflectors)
    - Wait for another router to change best route
- ♦ Transmitting many BGP messages
  - Latency for transferring the data

Delays very sensitive to router implementation decisions and network design!

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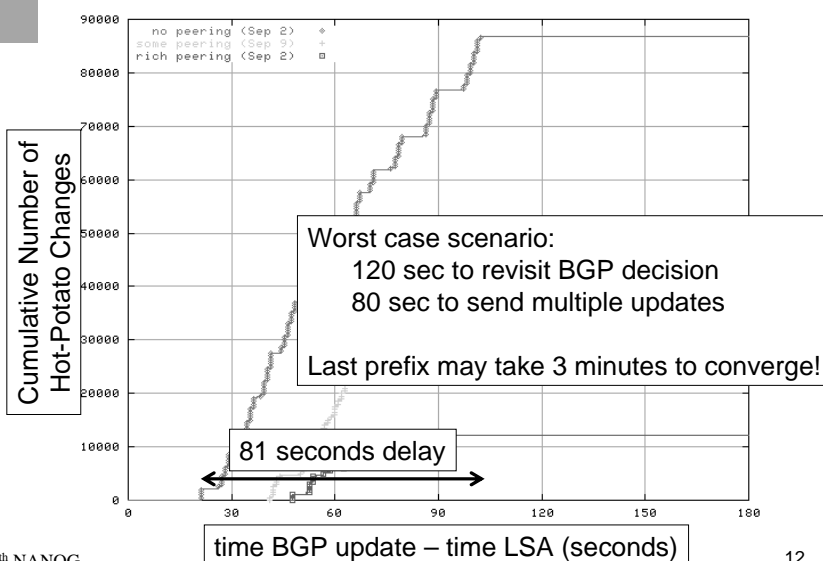
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## Delay for First BGP Update



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## Transferring Multiple Prefixes



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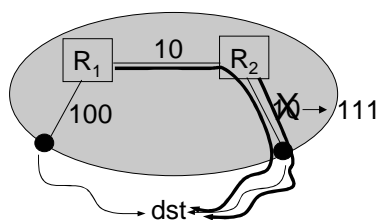
## Implications on Packet Forwarding

- ◆ Forwarding plane convergence
  - FIB update only when both IGP and BGP converge
- ◆ Measuring the performance impact
  - Do we need new measurement techniques?

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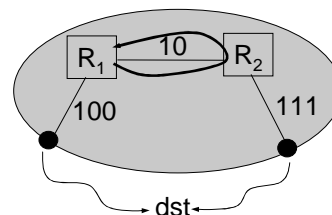
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## Forwarding Plane Convergence



- 1 - Scan process runs in R<sub>2</sub>
- 2 - R<sub>2</sub> starts using R<sub>1</sub> to reach dst
- 3 - R<sub>1</sub>'s scan process can take up to 60 seconds to run

Packets to dst may be caught in a loop for 60 seconds!

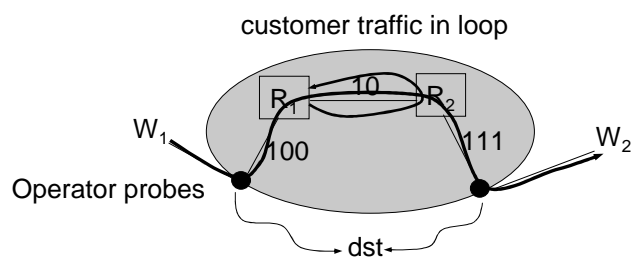


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## Detecting Loops is Hard

- ◆ Active probing of forwarding path
  - Probing machines just have one exit point



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## Avoiding Unnecessary Hot-Potato Changes

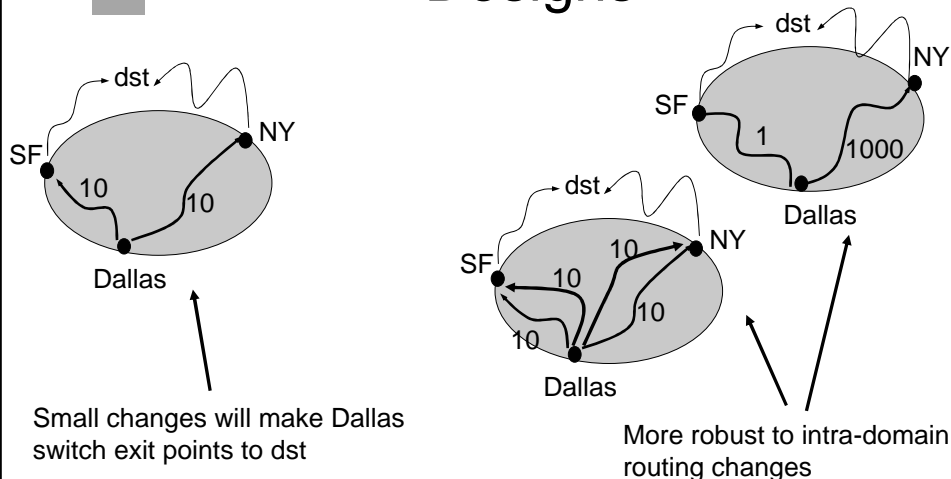
- ◆ Router connectivity to exit points
  - Likelihood of hot-potato routing changes
- ◆ Cost in/out of links during maintenance

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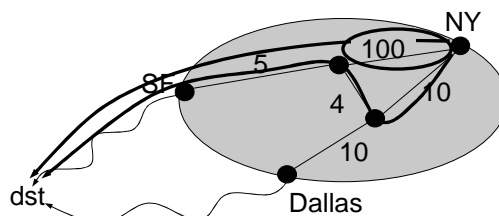
## Comparison of Network Designs



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## Careful Cost in/out Links



Traffic is more predictable  
Faster convergence  
Less impact on neighbors

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## Conclusions

- ♦ Hot potatoes can affect a lot of prefixes and have high convergence delays
- ♦ Vendors:
  - Event-driven implementations
  - Understand network-wide effect
- ♦ Operators:
  - Tuning timer
  - Network design
  - Maintenance practices



## Future Work

- ♦ Impact of the IGP-triggered BGP updates
  - Sudden shifts in the flow of traffic
  - Forwarding loops during convergence
  - Externally visible BGP routing changes